Evolution of testing techniques: From active testing to monitoring techniques

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- Show the evolution of active testing to monitoring (passive testing) techniques
- Explain the differences and complementarity of these techniques
- Present some representative examples

## A collaborative model of research

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- Our research model is based in:
  - Basic and applied research
  - Evaluation of results in real environments
  - Strong collaboration with industrial partners



## **Different Application Domains**











DETAILED WEB SERVICES PROCESS



Figure 1: The process flow of a Web service





- Testing: The process of executing software with the intent of finding and correcting faults
- Conformance testing: The process of checking if the implementation under test conforms the specification
  - Two techniques: active and passive testing (monitoring)
  - This presentation will focus mostly on monitoring, but there are many common objectives and challenges with active testing



- Usually called Model Based Testing (MBT)
- It is assumed that the tester controls the implementation. Control means: after sending an input and after receiving an output, the tester knows what is the next input to be send
- The tester can guide the implementation towards specific states
- Automatic test generation methods can be defined
- Usually a test case is a set of input sequences



- Passive testing consists in analyzing the traces recorded from the IUT and trying to find a fault by comparing these traces with either the complete specification or by verifying some specifics requirements (or properties) during normal runtime
- No interferences with the IUT
- It is also referred to as monitoring







## **Controllability issue in active testing**

- How to bring the finite state machine implementation into any given state at any given time during testing ?
  - Non trivial problem because of limited controllability of the finite state machine implementation
  - It may not be possible to put the finite state machine into the head state of the transition being tested without realizing several transitions



### **Observability issue in testing**

- How to verify that the finite state machine implementation is in a correct state after input/output exchange?
  - *State identification* problem. Difficult because of limited observability of the finite state machine implementation, it may not be possible to directly verify that the finite state machine is in the desired tail state after the transition has been fired

# Solutions to observability issue

To solve this problem different methods have been proposed:
DS (Distinguishing Sequence)
UIO (Unique Input/Output Sequence)
W (Distinction Set)

### **Unique Input/Output sequence** (UIO sequence)

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c/x

**S**1

Define an input sequence for each state such that the output sequence generated is unique to that state. Detects output and transfer faults.



#### **Transfer and output error detection**

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Test of (1): a/y a/x b/y Test of (2): a/y c/z b/y

Application du test of (1) to the implementation: a/y a/x b/z (transfer error) Application of test (2) to the implementation:

a/y c/x (output error)

# Limitations of active testing

- Non applicable when no direct access to the implementation under test
- Semi- controllable interfaces (component testing)
- Interferences on the behaviour of the implementation

## **Components Testing**

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#### Test in context, embedded testing:

- Tests focused on some components of the system, to avoid redundant tests
- O Interfaces semi-controllables
- In some cases it is not possible to apply active testing



## Why passive testing?

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- Conformance testing is essentially focused on verifying the conformity of a given implementation to its specification
  - It is based on the ability of a tester that stimulates the implementation under test and checks the correction of the answers provided by the implementation

#### Closely related to the controllability of the IUT

- In some cases this activity becomes difficult, in particular:
  - ★ if the tester has not a direct interface with the implementation
  - or when the implementation is built from components that have to run in their environment and cannot be shutdown or interrupted (for long time) in order to test them

# Controllability and observability issues in passive testing

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#### Controllability

• No **controllability** issue because no interaction with the implementation under test

#### Observability

• It is assumed that to perform passive testing it is necessary to observe the messages exchanges between modules.

• Passive testing is a Grey Box testing technique

#### Fault detection using passive testing

• It is possible to detect output faults

• It is possible to detect transfer faults under some hypothesis: to initialise the IUT in order to be sure that the implementation is in the initial state and then perform passive testing

#### **Invariant based passive testing approach**

- In this approach a set of properties are extracted from the specification or proposed by the protocol experts, and then the trace resulting from the implementation is analyzed to determine whether it validates this set of properties.
- These extracted set of properties are called invariants because they have to hold true at every moment.

#### **Invariant based passive testing approach**

- Definition: an invariant is a property that is always true.
- Two test steps:
  - Extraction of invariants from the specification or proposed by protocol experts
  - Application of invariants on execution event traces from implementation
- Solution: I/O invariants

## **Test by invariants: I/O invariants**

- An invariant is composed of two parts :
  - The test (an input or an output)
  - The preamble (I/O sequence)
- 3 kind of invariants :
  - o Output invariant (simple invariant)
  - Input invariant (obligation invariant)
  - Succession invariant (loop invariant)

## **Test by invariants : Simple (Output)** invariant

- Definition : invariant in which the test is an output
- Meaning : « immediatly after the sequence *préambule* there is always the expected output »
- Example :

 $(i_1 / o_1) (i_2 / o_2)$ (preambule in blue, expected output in red)

## Test by invariants : Obligation (Input) invariant

- Definition : invariant in which the test is an input
- Meaning : « immediatly before the sequence *preamble* there is always the input *test »*
- Example :

 $(i_1 / o_1) (i_2 / o_2)$ (preamble in blue, test in red)

#### **Test by invariants : succession invariant**

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- Definition : I/O invariant for complex properties (loops ...)
- Example :

the 3 invariants below build the property :
 « only the third i<sub>2</sub> is followed by o<sub>3</sub> »
 (i<sub>1</sub> / o<sub>1</sub>) (i<sub>2</sub> / o<sub>2</sub>)

 $(i_1 / o_1) (i_2 / o_2) (i_2 / o_2)$  $(i_1 / o_1) (i_2 / o_2) (i_2 / o_2) (i_2 / o_3)$ 

#### **Simple invariant**

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- A trace as *i*<sub>1</sub>/*O*<sub>1</sub>,..., *i*<sub>n-1</sub>/*O*<sub>n-1</sub>, *i*<sub>n</sub>/*O* is a simple invariant if each time that the trace *i*<sub>1</sub>/*O*<sub>1</sub>,..., *i*<sub>n-1</sub>/*O*<sub>n-1</sub> is observed, if we obtain the input *i*<sub>n</sub> then we necessarily get an output belonging to *O*, where *O* is included in the set of expected outputs.
- *i/o, \*, i'/O* means that if we detect the transition *i/o* then the first occurrence of the symbol *i*' is followed by an output belonging to the set *O*.
- \* replaces any sequence of symbols not containing the input symbol *i*' and ? replaces any input or output.





## **Run Time Verification**

- Approach proposed by researchers of verification (model checking) community
- Passive testing developed by the testing community
- EAGLE and RuleR tools proposed by Barringer and al. in 2004 and 2010 respectively, based on temporal logics and rewriting rules for properties description
- Others tools: Tracematches, [Avgustinov et al. 2007], J-LO [Bodden 2005]and LSC [Maoz and Harel 2006]



## How we see monitoring

- Monitoring the traces of a running system (e.g., traffic or message flows), **online or offline**.
- **Non-obtrusive** (i.e., execution traces are observed without interfering with the behaviour of the system).
- Analyzing collected data according to functional and non-functional requirements:
  - **Security properties** described in a formal specification (temporal logic , regular expressions, describing behaviour involving several events over time).
  - **Performance** to get real time visibility over the traffic statistics, KPI, delivered QoS, etc.
- Extended to perform **counter-measures**.











- INTER-TRUST project. Three-year project with many academic and industrial partners
- Security properties of services
- Detection of attacks using active and monitoring techniques

#### MOTIVATION

#### • Why testing ? (testing phase)

- Vulnerabilities can be introduced by AOP (Aspect Oriented Programming) used in Inter-trust
  - × Functional testing
  - × Check the respect of weaved security policies (aspects)
  - × Check the robustness of the target application
  - × Detect vulnerabilities
  - × Simulate attacks

#### • Why monitoring ? (testing & operation phases)

- Same as above
- + detecting context changes (context awareness) at runtime





- Generation of tests from IF model and test purposes Target: functional, security properties, attacks
- Execution relying on Selenium (Web interface)
- Detecting failures using MMT

Elections - Logged as Bob	ord) / access denied	F FSM Model
Licentinis Logged as Bob		
Vote Verification		iser oose ons
Question L1.1— Log Out Signed in as login1		
Q1: Which of the I Voting Receipt		Step 3 of 3
A1: Sports Question L12 Q2: Do you and The E-voting application has been specified as an extended finite state machine (IF language) Q2: Do you agree with the prop ST. The E-voting application has been specified as an extended finite state machine (IF language) Q2: Do you agree with the prop ST. The E-voting application has been specified as an extended finite state machine (IF language) Q2: Do you agree with the prop ST. The E-voting application has been specified as an extended finite state machine (IF language) Q2: Do you agree with the prop	nent. We recommend that you print it. [e12Eame/NmPI+uIkB azEHdrXas8CZyPXLvhVH/UhDog8Qznh 6UNyZYqNfaq8UCYvmpvem5sCQEHoJ8 Dk5BGA1rex8NJqLu5CYfV8JNn v4dWbFuYCDK2xkSRfAxv5v16J01jBge +w56Qe9dI61e19Rdk2eKv9 osal of having acces to a virtual newspaper in working hours?	f kT94w9mkWatChben9e3EEVCYnzDD+mK fSmf8l 21Vtdi6VVk/h+iinv 9KurHgKv8 tW This state presents the In this step the vote choices are displayed. The user has to fill the vote form. The step is the effective vote



#### **Test Case Generation**



### **Automatic Test Execution**





## Monitoring

#### • Two main uses:

- During the testing phase to complement the testing tools and provide a verdict
- During the operation phase to monitor security and application context
- Relies on data collected at different levels
  Network (ex. CAM messages)
  Application internal events (notification module)
  System status (CPU and memory usage)

#### **Montimage Monitoring Tool**



#### **Analysis and Failure Detection**

- Evoting test case Advanced authentication option
  - Example of property: Only authenticated voters can cast their votes



## MMT Analysis Dashboard-Security Property

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#### Security Dashboard



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## MMT Analysis Dashboard-Attack Detection



Security rule	Description
<b>R2=</b> Permission(Org1, voter, access, L1, intranet_ context)	A voter connecting from the intranet system is permitted to access to the first type of election(named L1)
R3=Permission(Org1, voter, access, L2, intranet_ context)	A voter connecting from the intranet system is permitted to access to the second type of election(named L2)
<b>R4=</b> Permission(Org1, voter, access, L3, intranet_ context)	A voter connecting from the intranet system is permitted to access to the second type of election(named L3)
<b>R5=</b> Permission(Org1, voter, access, L5, desktop_ context)	A voter connecting from a specific desktop is permitted to access to the second type of election(named L4)
<b>R6=</b> Permission(Org1, voter, access, L4, desktop_ context)	A voter connecting from a specific desktop is permitted to access to athe second type of election(named L5)
<b>R7=</b> Obligation(Org1, server, authorize, Election_Lists, violation_context)	The server is authorized to check the elections lists when a violation is detected
<b>R</b> 8=Obligation(Org1, voter, authenticate, L1, normal_context)	The voter should be authenticated to access to the first type of election (L1) when the normal context is activated (anonymous context is not activated).
<b>R9=</b> Obligation (Org1, voter, encrypt, L1_vote, violation_context)	The vote device should encrypt the data related to any election with the type L1 when a violation is detected.
<b>R10=</b> Obligation(Org1, voter, authenticate, L3, violation_context)	The voter should be authenticated to access to the third type of election (L3) when a violation is detected.
<b>R</b> 11=Obligation(Org1, voter, Advanced_Authenticate, L4, normal_context, violation_context)	The voter should be authenticated based on advanced methods to access to the third type of

#### Summary

- Model based test generation for security purposes (TestGen-IF)
- Correlation of data from different sources (Network, application, system)
- Detection of attacks and failures at runtime reaction
- Brings dynamicity to system by adapting to different contexts

## **Related Works**

- Monitoring of routing protocols for ad hoc (OLSR protocol) and mesh networks networks based on a dsitributed approach (Batman protocol) (Telecom Sud Paris)
- Monitoring for secure interoperability Application to a multi-source information system (Telecom Sud Paris)
- Monitoring with time constraints (C. Andrés, M. Nuñez and Mercedes Merayo)
- Monitoring with AsynchronousCommunications (M. Nuñez and R. Hierons)
- Other works by (T. Jeron and H. Marchand, A. Ulrich and A. Petrenko)

#### Conclusions

- It is now easier to support active testing and monitoring and integrate it with other development activities
- Modeling technology has matured (using FSMs, EFSMs, different UML profiles (SysML), temporal logics)
- Much research and innovation is still required and it should involve collaborations between research and industry

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